

On a special structured matrix problem

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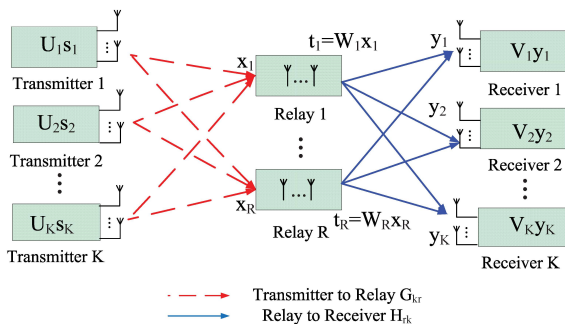
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- 1 Introduction
- 2 TSTINR approximation model
- 3 Multiple stream TSTINR model
- 4 Simulation results
- 5 Conclusion

Applications



- multiple user pairs, multiple relays, multiple antennas

Transmission process

Time slot 1: transmitters \rightarrow relays

Time slot 2: relays \rightarrow receivers

$$\tilde{\mathbf{y}}_k = \underbrace{\mathbf{V}_k^H \mathbf{T}_{kk} \mathbf{s}_k}_{\text{desired signal}} + \underbrace{\sum_{q \in \mathcal{K}, q \neq k} \mathbf{V}_k^H \mathbf{T}_{kq} \mathbf{s}_q}_{\text{interference}} + \underbrace{\sum_{r \in \mathcal{R}} \mathbf{V}_k^H \mathbf{H}_{kr} \mathbf{W}_r \mathbf{n}_r + \mathbf{V}_k^H \mathbf{z}_k}_{\text{noise}} \quad (1.1)$$

where $\mathbf{T}_{kq} = \sum_{r \in \mathcal{R}} \mathbf{H}_{kr} \mathbf{W}_r \mathbf{G}_{rq} \mathbf{U}_q$.

- Optimize:
 - precoding matrices $\{\mathbf{U}\}$
 - decoding matrices $\{\mathbf{V}\}$
 - relay AF matrix $\{\mathbf{W}\}$

Sum rate maximization

- Aim: **maximize** system sum rate under certain power constraints

$$R_{\text{sum}} = \frac{1}{2} \sum_{k \in \mathcal{K}} \log_2 \det(\mathbf{I}_{N_k} + \mathbf{F}_k^{-1} \mathbf{T}_{kk} \mathbf{T}_{kk}^H), \quad (1.2)$$

where $\mathbf{F}_k = \sum_{q \neq k, q \in \mathcal{K}} \mathbf{T}_{kq} \mathbf{T}_{kq}^H + \sum_{r \in \mathcal{R}} \sigma_r^2 \mathbf{H}_{kr} \mathbf{W}_r (\mathbf{H}_{kr} \mathbf{W}_r)^H + \mu_k^2 \mathbf{I}_{N_k}$.

- ☹ Difficult to optimize directly
- ☺ Approximation model to approximate sum rate maximization

- Total Signal to Total Interference plus Noise Ratio (**TSTINR**) definition:

$$\text{TSTINR} = \frac{P^S}{P^I + P^N} = \frac{\sum_{k \in \mathcal{K}} P_k^S}{\sum_{k \in \mathcal{K}} (P_k^I + P_k^N)}.$$

Theorem 2.1

For any precoder $\mathbf{U}_k, k \in \mathcal{K}$, relay beamforming matrix $\mathbf{W}_r, r \in \mathcal{R}$ and any decoder \mathbf{V}_k satisfying $\mathbf{V}_k^H \mathbf{V}_k = \mathbf{I}_{d_k}, k \in \mathcal{K}$, we have

$$\log_2[1 + \text{TSTINR}(\{\mathbf{U}\}, \{\mathbf{V}\}, \{\mathbf{W}\})] \leq R_{\text{sum}}(\{\mathbf{U}\}, \{\mathbf{W}\}).$$

- With individual user and relay transmit power constraints:

$$\begin{aligned} \max_{\substack{\{\mathbf{U}\}, \{\mathbf{V}\}, \\ \{\mathbf{W}\}}} \quad & \text{TSTINR} = \frac{\sum_{k \in \mathcal{K}} P_k^S}{\sum_{k \in \mathcal{K}} (P_k^I + P_k^N)} \\ \text{s.t.} \quad & \mathbf{V}_k^H \mathbf{V}_k = \mathbf{I}_{d_k}, \|\mathbf{U}_k\|_F^2 \leq p_k^U, k \in \mathcal{K}, \\ & \sum_{k \in \mathcal{K}} \|\mathbf{W}_r \mathbf{G}_{rk} \mathbf{U}_k\|_F^2 + \sigma_r^2 \|\mathbf{W}_r\|_F^2 \leq p_r^R, r \in \mathcal{R}. \end{aligned} \quad (2.1)$$

- Problem reformulation:

$$\begin{aligned} \min_{\substack{\{\mathbf{U}\}, \{\mathbf{V}\}, \\ \{\mathbf{W}\}}} \quad & f(\{\mathbf{U}\}, \{\mathbf{V}\}, \{\mathbf{W}\}; C) = C(P^I + P^N) - P^S \\ \text{s.t.} \quad & \dots \end{aligned} \quad (2.2)$$

- Update strategy of the parameter C :

$$C = \frac{P^S(\{\mathbf{U}\}, \{\mathbf{V}\}, \{\mathbf{W}\})}{P^I(\{\mathbf{U}\}, \{\mathbf{V}\}, \{\mathbf{W}\}) + P^N(\{\mathbf{U}\}, \{\mathbf{V}\}, \{\mathbf{W}\})}. \quad (2.3)$$

Theorem 2.2

If the objective function of (2.2) has sufficient reduction in each iteration and C is updated as (2.3), then the objective function of (2.1), TSTINR, is monotonically increasing. Any KKT point of (2.2) is also a KKT point of (2.1).

Alternating direction algorithm

- Decoding subproblem:

$$\begin{aligned} \min_{\mathbf{X} \in \mathbb{C}^{N_k \times d_k}} \quad & \text{tr}(\mathbf{X}^H \mathbf{A}_0 \mathbf{X}) \\ \text{s.t.} \quad & \mathbf{X}^H \mathbf{X} = \mathbf{I}_{d_k} \end{aligned} \quad (2.4)$$

- Eigenvalue problem

- Equivalent subproblem for relay AF matrix \mathbf{W}_r :

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{C}^{L_r^2 \times 1}} \quad & \tilde{f}(\mathbf{x}) = \mathbf{x}^H \mathbf{B}_1 \mathbf{x} + \mathbf{b}^H \mathbf{x} + \mathbf{x}^H \mathbf{b} \\ \text{s.t.} \quad & \mathbf{x}^H \mathbf{B}_2 \mathbf{x} \leq \eta_1, \end{aligned} \quad (2.5)$$

where $\mathbf{B}_2 > 0$.

- Trust region subproblem

Alternating iteration algorithm

- Equivalent subproblem for precoder \mathbf{U}_k :

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^{n \times 1}} \quad & f(\mathbf{x}) = \mathbf{x}^T \mathbf{Q}_0 \mathbf{x} + 2\mathbf{g}^T \mathbf{x} \\ \text{s.t.} \quad & C_i(\mathbf{x}) = \mathbf{x}^T \mathbf{Q}_i \mathbf{x} - 1 \leq 0, i = 1, \dots, m, \end{aligned} \quad (2.6)$$

- \mathbf{Q}_0 indefinite; $\mathbf{Q}_i \geq 0, i = 1, \dots, m$
- Nonconvex Quadratic Constrained Quadratic Programming (QCQP)

Idea of feasible shrinkage method

- Initial point:

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^{n \times 1}} \quad & f(\mathbf{x}) = \mathbf{x}^T \mathbf{Q}_0 \mathbf{x} + 2\mathbf{g}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{x}^T \sum_{i=1}^m \mathbf{Q}_i \mathbf{x} \leq 1 \end{aligned} \quad (2.7)$$

- The optimal solution of (2.7) is a feasible solution of (2.6).

For any \mathbf{x} as the feasible solution of (2.7),

$$\mathbf{x}^T \mathbf{Q}_i \mathbf{x} \leq \mathbf{x}^T \sum_{i=1}^m \mathbf{Q}_i \mathbf{x} \leq 1, i = 1, \dots, m.$$

Idea of feasible shrinkage method

- Subproblem in each iteration (with the iterative point \mathbf{x}_0):

$$\begin{aligned} \min_{\mathbf{d} \in \mathbb{C}^{n \times 1}} \quad & \mathbf{d}^H \mathbf{Q}_0 \mathbf{d} + \bar{\mathbf{g}}^H \mathbf{d} + \mathbf{d}^H \bar{\mathbf{g}} \\ \text{s.t.} \quad & \mathbf{d}^H \bar{\mathbf{Q}} \mathbf{d} \leq 1, \end{aligned} \quad (2.8)$$

where $\bar{\mathbf{g}} = \mathbf{g} + \mathbf{Q}_0 \mathbf{x}_0$, $\bar{\mathbf{Q}} = \sum_{i=1}^m \rho_i (\mathbf{Q}_i + \rho_i \mathbf{Q}_i \mathbf{x}_0 \mathbf{x}_0^H \mathbf{Q}_i)$ and

$$\rho_i = \frac{2}{1 - \mathbf{x}_0^H \mathbf{Q}_i \mathbf{x}_0}, i = 1, \dots, m.$$

- ✘ Subproblem idea: weighted quadratic approximation

Feasible shrinkage method

Lemma 2.1

Let $\mathbf{A} \in \mathbb{R}^{n \times n}$ be a positive semi-definite symmetric matrix and $\mathbf{y} \in \mathbb{R}^{n \times 1}$ is a vector, such that $\mathbf{y}^T \mathbf{A} \mathbf{y} < 1$, then

$$\mathcal{B} = \{\mathbf{x} | (\mathbf{x} - \mathbf{y})^T (\mathbf{A} + \rho \mathbf{A} \mathbf{y} \mathbf{y}^T \mathbf{A}) (\mathbf{x} - \mathbf{y}) \leq \frac{1}{\rho}\}$$

is the subset of $\mathcal{A} = \{\mathbf{x} | \mathbf{x}^T \mathbf{A} \mathbf{x} \leq 1\}$, where $\rho = \frac{2}{1 - \mathbf{y}^T \mathbf{A} \mathbf{y}}$.

Theorem 2.3

If \mathbf{x}_0 is an interior feasible point of (2.6) and \mathbf{d}_0 is the optimal solution of (2.8), then $\mathbf{x}_0 + \mathbf{d}_0$ is feasible for problem (2.6).

Algorithm 1 (Feasible shrinkage algorithm)

- Step 1** Initialization: given stopping parameters ϵ_0 and ϵ_1 ; solve (2.7) and set its optimal solution as \mathbf{x}_0 .
- Step 2** With the feasible iterative point \mathbf{x}_0 , solve subproblem (2.8).
- Step 3** Update the iterative point with \mathbf{d}_0 , the solution of (2.8), as $\mathbf{x}_0 := \mathbf{x}_0 + \mathbf{d}_0$.
- Step 4** If $\|\mathbf{d}_0\|_2 < \epsilon_1$, or there exists i such that $1 - \mathbf{x}_0^T \mathbf{Q}_i \mathbf{x}_0 < \epsilon_0$, stop and output \mathbf{x}_0 . Otherwise go back to Step 2.

- feasible shrinkage method + nonconvex SQP
- interior point \rightarrow boundary point
- a good feasible solution \rightarrow KKT point

Algorithm 2 (Algorithm for TSTINR model)

Step 1 Set initial value of $\mathbf{U}_k, k \in \mathcal{K}$ and $\mathbf{W}_r, r \in \mathcal{R}$. $C = 1$.

Step 2 Update decoder \mathbf{V}_k by solving (2.4), $k \in \mathcal{K}$.

Step 3 Update relay beamforming matrix \mathbf{W}_r by solving (2.5), $r \in \mathcal{R}$.

Step 4 Update precoder \mathbf{U}_k by solving (2.6), $k \in \mathcal{K}$.

Step 5 Update C as $C := \frac{P^S}{P^I + P^N}$. Go back to Step 2. Iterate until convergence.

- ☺ The objective function TSTINR converges
- ☹ No guarantee to converge to a KKT point

Analysis of precoding subproblem

- Precoding subproblem:

$$\begin{aligned} \min_{\mathbf{X} \in \mathbb{C}^{M_k \times d_k}} \quad & \text{tr}(\mathbf{X}^H \mathbf{D}_0 \mathbf{X}) \\ \text{s.t.} \quad & \text{tr}(\mathbf{X}^H \mathbf{D}_r \mathbf{X}) \leq \eta_r, r \in \mathcal{R}. \end{aligned} \quad (3.1)$$

Here $\mathbf{D}_r \geq 0, \eta_r \geq 0, r \in \mathcal{R}$

Theorem 3.1

The precoding subproblem (3.1) always has a rank one optimal solution.

- Observation: rank one precoding matrices
linearly dependent columns

Multiple stream model

$$\begin{aligned} \max_{\substack{\{\mathbf{U}\}, \{\mathbf{V}\}, \\ \{\mathbf{W}\}}} \quad & \text{TSTINR} = \frac{P^S}{P^I + P^N} = \frac{\sum_{k \in \mathcal{K}} P_k^S}{\sum_{k \in \mathcal{K}} (P_k^I + P_k^N)} \quad (3.2) \\ \text{s.t.} \quad & \mathbf{U}_k^H \mathbf{U}_k = \frac{p_k^U}{d_k} \mathbf{I}_{d_k}, \leftarrow \text{Different from (2.1)} \\ & \mathbf{V}_k^H \mathbf{V}_k = \mathbf{I}_{d_k}, k \in \mathcal{K}, \\ & \sum_{k \in \mathcal{K}} \|\mathbf{W}_r \mathbf{G}_{rk} \mathbf{U}_k\|_F^2 + \sigma_r^2 \|\mathbf{W}_r\|_F^2 \leq p_r^R, r \in \mathcal{R} \end{aligned}$$

- Precoding subproblem:

$$\min_{\mathbf{X} \in \mathbb{C}^{M_k \times d_k}} \quad \text{tr}(\mathbf{X}^H \mathbf{D}_0 \mathbf{X}) \quad (3.3a)$$

$$\text{s.t.} \quad \mathbf{X}^H \mathbf{X} = \rho \mathbf{I}_{d_k}, \quad (3.3b)$$

$$\text{tr}(\mathbf{X}^H \mathbf{D}_r \mathbf{X}) \leq \eta_r, r \in \mathcal{R} \quad (3.3c)$$

- \mathbf{D}_0 indefinite; $\mathbf{D}_r \geq 0, r \in \mathcal{R}$
- Nonconvex QCQP with orthogonality constraint

Dual problem

- Dual Problem of (3.3):

$$\begin{aligned} \min_{\theta_r, r \in \mathcal{R}} \quad & h(\theta_1, \theta_2, \dots, \theta_R) \\ & = \sum_{r \in \mathcal{R}} \theta_r \eta_r - \min_{\mathbf{X}^H \mathbf{X} = \rho \mathbf{I}_{d_k}} \text{tr}[\mathbf{X}^H (\mathbf{D}_0 + \sum_{r \in \mathcal{R}} \theta_r \mathbf{D}_r) \mathbf{X}] \\ \text{s.t.} \quad & \theta_r \geq 0, r \in \mathcal{R}. \end{aligned} \tag{3.4}$$

Lagrangian multiplier: $\boldsymbol{\theta} = (\theta_1, \theta_2, \dots, \theta_R)^T$

- Suppose $\rho = 1$
- ✓ Projected subgradient method

Projected subgradient method

- Subgradient of $h(\boldsymbol{\theta})$: $\mathbf{y} = (y_1, y_2, \dots, y_R)^T$,

$$y_r = \frac{\partial h}{\partial \theta_r} = \eta_r - \text{tr}(\mathbf{X}^H \mathbf{D}_r \mathbf{X})$$

- \mathbf{X} : eigenvectors of the smallest d_k eigenvalue of $\mathbf{D}_0 + \sum_{r \in \mathcal{R}} \theta_r \mathbf{D}_r$

Ref Rudisill, "Derivatives of Eigenvalues and Eigenvectors for a General Matrix," AIAA J., 1974

Algorithm 3 (Algorithm for subproblem (3.3))

Step 1 Randomly set initial point $\boldsymbol{\theta} \geq \mathbf{0}$. Set the stopping criterion $\epsilon \geq 0$.
 $j = 1$.

Step 2 Update the iterative point: $\boldsymbol{\theta}^{j+1} = (\boldsymbol{\theta}^j - \frac{1}{j} \mathbf{y}^j)_+$. $j := j + 1$.

Step 3 If $\|\boldsymbol{\theta}^{j+1} - \boldsymbol{\theta}^j\|_2 < \epsilon$, stop and output $\boldsymbol{\theta}^* = \boldsymbol{\theta}^j$ as the solution of (3.4) and $\mathbf{X}^* = \mathbf{v}_{\min}^{d_k}(\mathbf{D}_0 + \sum_{r \in \mathcal{R}} \theta_r^* \mathbf{D}_r)$ as the solution of (3.3); otherwise, go back to Step 2.

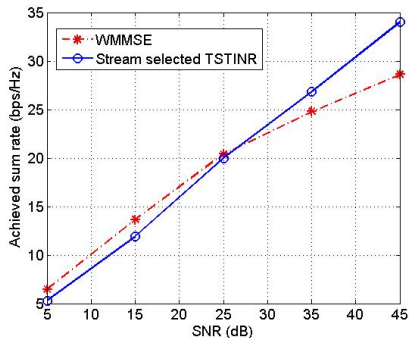
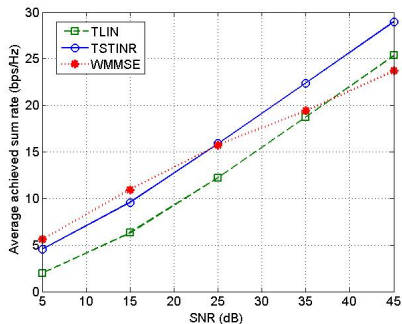
Theorem 3.2

If θ^ is the optimal solution of the dual problem (3.4), then \mathbf{X}^* is a feasible point of (3.3). Furthermore, if $h(\theta)$ is smooth at $\theta = \theta^*$, then \mathbf{X}^* is the optimal solution of (3.3).*

Simulation settings

- Each element of \mathbf{G}_{rk} and $\mathbf{H}_{kr}, k \in \mathcal{K}, r \in \mathcal{R}: \mathcal{CN}(0, 1)$
- $\sigma_1^2 = \sigma_2^2 = \sigma^2 = 1, \text{SNR} = \frac{p_0^U}{\sigma^2} = \frac{p_0^R}{\sigma^2}$.
- Initialization: randomly generated initial points;
penalty parameter $C = 1$.
- FS: $\epsilon_0 = 10^{-4}, \epsilon_0 = 10^{-8}$; SQP: $\epsilon = 10^{-8}, u = 1, \eta = 0.1$.
- 100 random realization of different channel coefficients for each plotted point

Comparison of different approaches



● Left: single stream; Right: multiple stream

✓ TSTINR achieves highest sum rate under medium to high SNR

Ref Truong, Sartori and Heath, "Cooperative Algorithms for MIMO Amplify-and-Forward Relay Networks," *IEEE Trans. Signal Process.*, 2013.

Comparison of different algorithms

SNR value	5dB	15dB	25dB	35dB	45dB
H-Sum rate (bps/Hz)	8.4579	14.1468	18.8045	23.4674	28.1119
C-Sum rate (bps/Hz)	8.4580	14.1468	18.8201	23.5132	28.1319
H-Time (s)	5.4138	5.5940	6.0208	5.9795	6.6703
C-Time (s)	106.3531	105.7086	106.3560	106.2938	107.1608

- convex subproblem: FS+ nonconvex SQP (H) **vs** CVX (C)

✓ Significantly less computing time with almost the same result

For the sum rate maximization problem in wireless communications:

- Propose efficient algorithms for the TSTINR model
(Algorithm for fractional optimization, special structured QCQP)
- Propose the multiple stream model and an efficient algorithm
(QCQP with orthogonal constraint)

Thank you for listening!